

Lesson 15: Boolean Representation of Ladder Diagrams

ET 438B Sequential Control and Data Acquisition
Department of Technology

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Learning Objectives

After this presentation you will be able to:

- Realize logic functions as ladder diagram rungs
- Follow the logic of a multi-rung ladder diagram
- Represent ladder rungs as Boolean gates
- Design combinational sequential controllers using Boolean equations

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Ladder Diagram Example

A manual mixing operation is to be automated using sequential process control methods. The process composed of three steps:

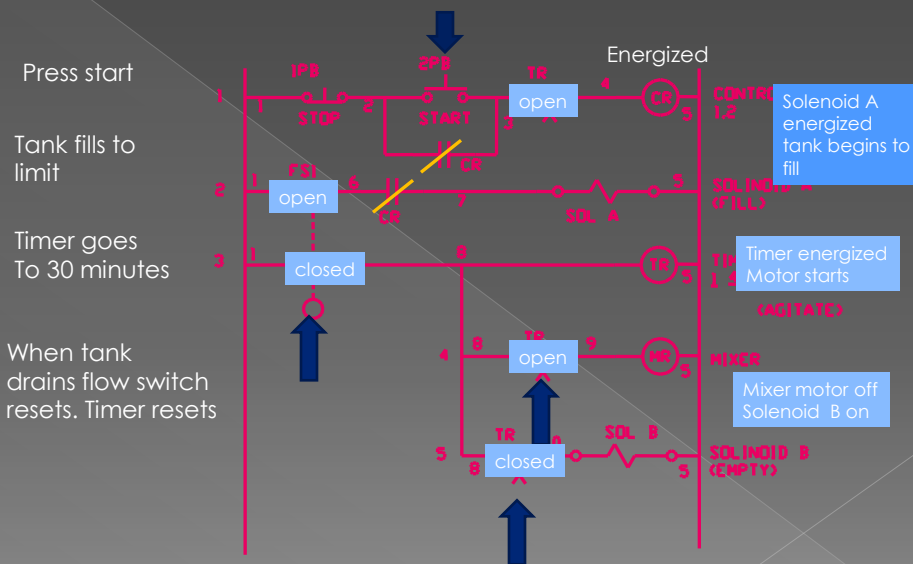
- filling a tank to a predetermined level
- agitating the liquid for 30 minutes
- draining the tank for use in another part of process

Does the ladder logic schematic that follows perform this function correctly?

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Ladder Diagram Example





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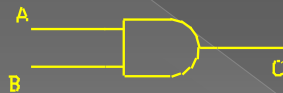
Combinational and Sequential Logic with Relays and Contacts

Let contact state represent a logical value

 = Logic 0 Called Form A Contact

 = Logic 1 Called Form B Contact

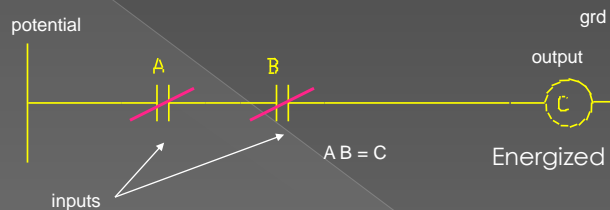
Implement AND gate



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Combinational and Sequential Logic with Relays and Contacts



Conditions A **AND** B must be present to energize output C

Note: all contacts are considered instantaneous and not held unless modified

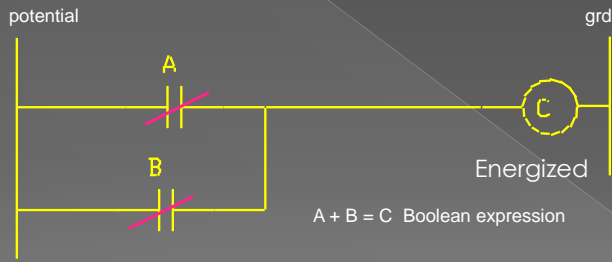
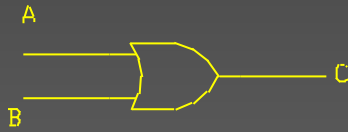
With electromechanical relays fan-in and fan-out limited by number of contacts in relays

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More Logic Functions

OR
Function



Either A **OR** B will cause coil C to be energized
Contacts A, B represent conditions or states in
the sequential process

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More Logic Functions

NOT Function



Boolean Expression
 $B = \overline{A}$



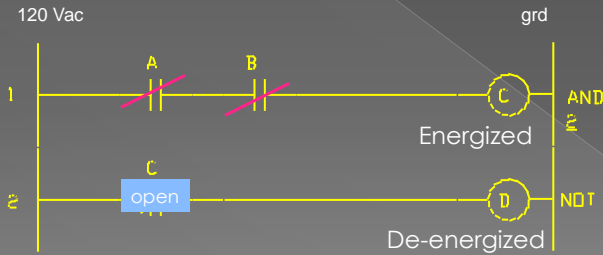
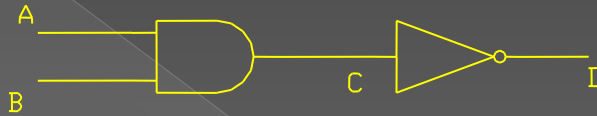
Contact of opposite state creates inversion

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Constructing Other Logic Functions

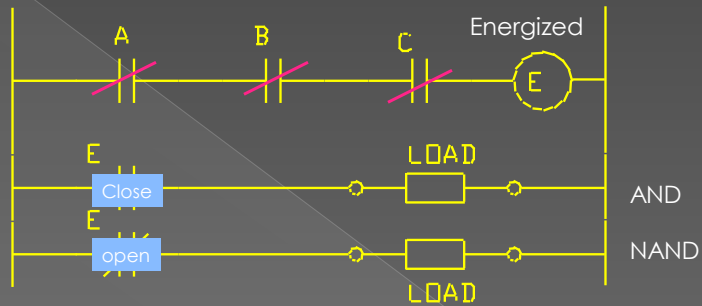
Combine the **AND** function with the **NOT** function to get a **NAND** operation.



Rung 1 implements the AND function
Rung 2 implements the NOT function

Any contact associated with coil D will change state like a NAND TTL gate.

Multiple Input AND/NAND

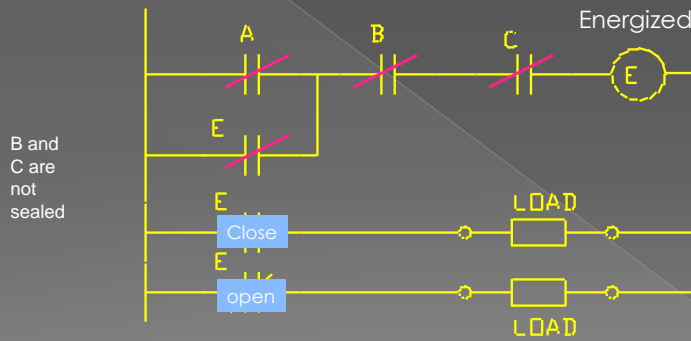


$$ABC = E \text{ and } \overline{ABC} = E$$

Can add a memory action to the above by including a feedback from the output coil to the inputs

Memory Action AND/NAND

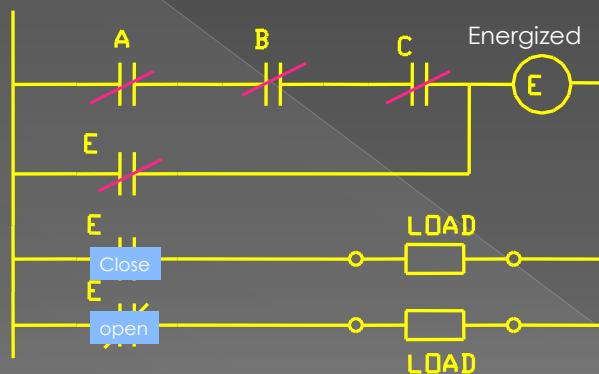
Can add a memory action to the above by including a feedback from the output coil to the inputs



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All Inputs Latched AND/NAND



The output can not change unless the circuit is de-energized.

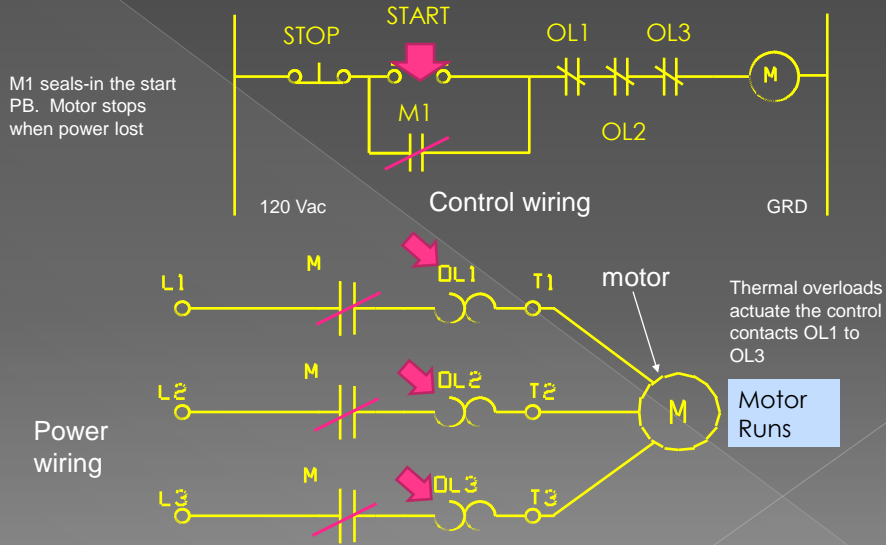
Contact E in rung 2 is a feedback from the output that makes circuit ignore state changes of A, B and C after the condition A B C is detected.

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Motor Control Example

Three-wire control- used for manual and automatic motor starting.



M1 seals-in the start PB. Motor stops when power lost

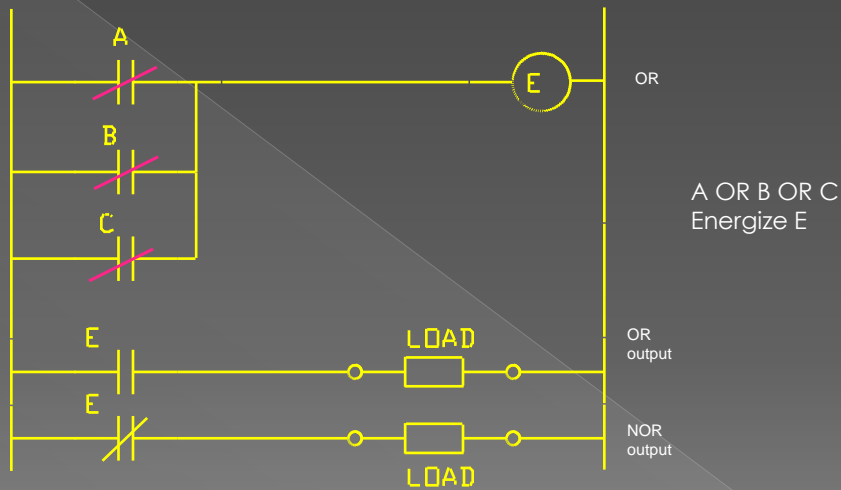
Thermal overloads actuate the control contacts OL1 to OL3

Motor Runs

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Multiple Input OR/NOR Function



Outputs $A+B+C = E$
 $A+B+C = E$

Notice that Relay logic is similar to TTL. Can use Truth tables and Boolean expressions to do designs

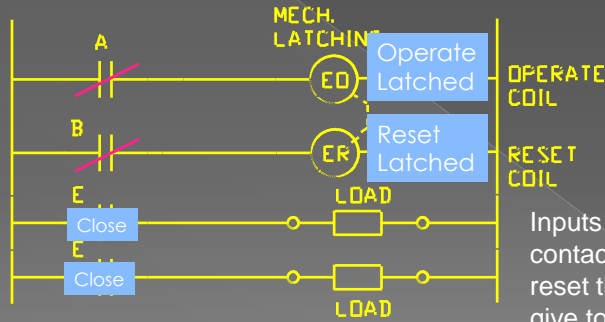
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Ladder Logic Memory Elements

Mechanically latched relay - maintains state even when power removed.
Has two coils (operate, reset)

Typical wiring



Inputs A and B set the output contacts E and reset then respectively. This give toggle action that "remembers" the last input state even when power is removed

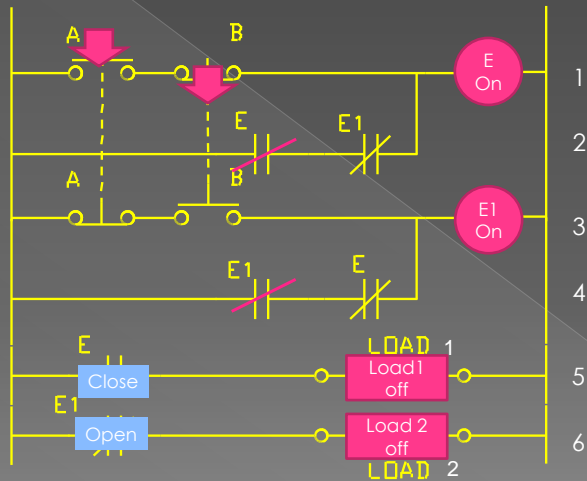
Typical Applications

Reversing Motor starters. Reclose Relay Cut-out

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Off-Return Memory



Energize and re-energize circuit - Load 2 on
No continuity in rungs 1-4
Continuity in rung 6

Press A: continuity rung 2 Both loads on

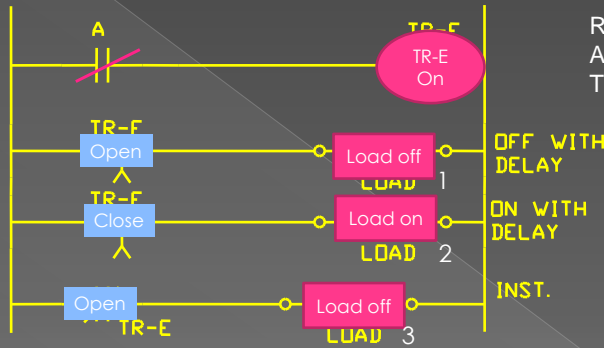
Press B: continuity rung 4 Both loads off

Remember all contacts are drawn with the coils de-energized

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Timer Sub-Circuits



Rung 1: when input A is energized timer TR-E starts

Schematic indicates that this is a on-delay timer. After defined interval TR-E in rung 2 opens and TR-E in rung 3 closes

Load 1 is deactivated after time delay
Load 2 is activated after time delay

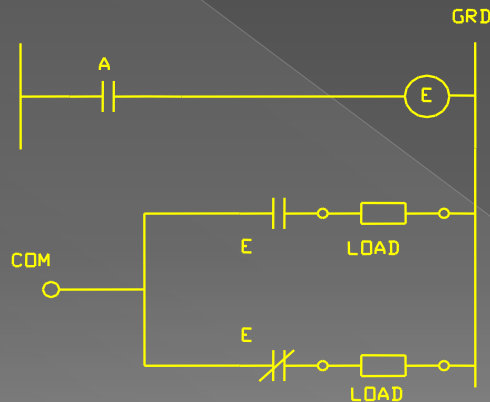
Load 3 is instantaneously deactivated by TR-E

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Form "C" Contact

Loads are toggled between a common point



Typical "Form C" contacts include both a NO and NC contact arrangement.

Used in some sensors for more flexibility

Contact A creates a remote control toggle switch

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Designing Sequential Control Systems

Combinational Systems

- Use true tables, Boolean Algebra
- Multiple inputs and/or outputs
- Sum of Products or product of sums Boolean Implementations
- Reduce to minimum implementation

Sequential Systems

- Follow steps, transition from one step to another.
- Use state transition diagrams or tables with Boolean Algebra
- State Machine implemented in software or hardware
- Decisions made base on current condition of system and input information

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Review of Logic Gates and Boolean Algebra

Boolean Variables False =0 True =1

Boolean Operators

AND



$$X = A \cdot B$$

A	B	X
0	0	0
0	1	0
1	0	0
1	1	1

OR



$$X = A + B$$

A	B	X
0	0	0
0	1	1
1	0	1
1	1	1

NOT



$$X = \bar{A}$$

A	X
0	1
1	0

EOR=XOR

NAND



$$X = \overline{A \cdot B}$$

A	B	X
0	0	1
0	1	1
1	0	1
1	1	0

NOR



$$X = \overline{A + B}$$

A	B	X
0	0	1
0	1	0
1	0	0
1	1	0

EOR



$$X = A \oplus B$$

A	B	X
0	0	0
0	1	1
1	0	1
1	1	0

Alternate Implementation

$$X = \bar{A}\bar{B} + \bar{A}B$$

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Review of Logic Gates and Boolean Algebra

Axioms of Boolean Algebra

Idempotent	Associative	Distributive	
$A + A = A$	$A + (B + C) = (A + B) + C$	$A + (B \cdot C) = (A + B) \cdot (A + C)$	
$A \cdot A = A$	$(A \cdot B) \cdot C = A \cdot (B \cdot C)$	$A \cdot (B + C) = (A \cdot B) + (A \cdot C)$	
Identity	Complement	DeMorgan's Theorem	Absorption
$A + 0 = A$	$A + \bar{A} = 1$	$\overline{A + B} = \bar{A} \cdot \bar{B}$	$A + \bar{A} \cdot B = A + B$
$A + 1 = 1$	$A \cdot \bar{A} = 0$	$\overline{A \cdot B} = \bar{A} + \bar{B}$	$A + A \cdot B = A$
$A \cdot 0 = 0$	$\overline{\bar{A}} = A$		Order of Operations 1. NOT 2. AND 3. OR
$A \cdot 1 = A$	$\bar{\bar{1}} = 0$		

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Review of Logic Gates and Boolean Algebra

Example: Simplify the following expression using the axioms of Boolean Algebra.

$$X = \overline{A + B \cdot C} + A(B + \bar{C})$$

$$X = \overline{A} + \overline{B \cdot C} + A(B + \bar{C})$$

Apply DeMorgan's Theorem to first term

$$\bar{A} + \overline{B \cdot C} = \bar{A} + (\bar{B} + \bar{C})$$

$$X = \bar{A} + \bar{B} + \bar{C} + A(B + \bar{C})$$

$$X = \bar{A} + \bar{B} + \bar{C} + A(B + \bar{C})$$

$$X = \bar{A} \cdot \bar{B} + \bar{A} \cdot \bar{C} + A \cdot B + A \cdot \bar{C}$$

Collect common terms and factor

$$\bar{C}(\bar{A} + A) = A \cdot \bar{C} + \bar{A} \cdot \bar{C}$$

Add
Parentheses

Apply
DeMorgan's
Here

$$\overline{B \cdot C} = \bar{B} + \bar{C}$$

Expand
Expressions

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Review of Logic Gates and Boolean Algebra

Example Continued

$$X = \bar{A} \cdot \bar{B} + A \cdot B + \bar{C} (A + \bar{A})$$

Use
Complement
Axiom

$$A + \bar{A} = 1$$

$$X = \bar{A} \cdot \bar{B} + A \cdot B + \bar{C} \cdot 1$$

Use Identity
Axiom

$$\bar{C} \cdot 1 = \bar{C}$$

$$X = \bar{A} \cdot \bar{B} + A \cdot B + \bar{C}$$

Simplified Expression

Logic Design

- 1.) Obtain description of process
- 2.) Define control action
- 3.) Define Inputs and Outputs
- 4.) Develop Truth Table or Boolean Equation of Process

Process control description

A heating oven with two bays can heat one ingot in each bay. When the heater is on it provides enough heat for two ingots. If only one ingot is present, the oven may overheat so a fan is used to cool the oven when it exceeds a set temperature.

Control Action

When only one ingot is in the oven and the temperature exceeds the setpoint, turn on the fan

Logic Design

Define I/O variables Inputs: B1 = bay1 ingot present
 B2 = bay2 ingot present
 T = temperature sensor

Create Truth Table

Output: F= fan start

T	B2	B1	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

If there is no over temperature
 don't start the fan

Over temperature in empty oven: safety fan start

Start fan in lightly load ovens with over temp.

Over temperature in full oven: safety fan start

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Logic Design

Select elements from truth table in SOP (sum-of-products) form then simplify.

T	B2	B1	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

$$F = T \cdot \bar{B1} \cdot \bar{B2} + T \cdot B1 \cdot \bar{B2} + T \cdot \bar{B1} \cdot B2 + T \cdot B1 \cdot B2$$

$$F = T \cdot (\bar{B1} \cdot \bar{B2} + B1 \cdot \bar{B2} + \bar{B1} \cdot B2 + B1 \cdot B2)$$

$$F = T \cdot (\bar{B2} \cdot (\bar{B1} + B1) + B2 \cdot (\bar{B1} + B1))$$

$$F = T \cdot (\bar{B2} + B2)$$

$$F = T$$

Requires only Temp
 control

Ignore unloaded and full load cases and try again

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Logic Design

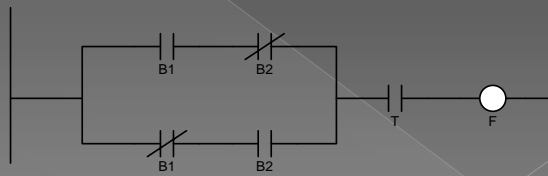
Revised Truth Table

T	B2	B1	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0

$$F = T \cdot B1 \cdot \overline{B2} + T \cdot \overline{B1} \cdot B2$$

$$F = T (B1 \cdot \overline{B2} + \overline{B1} \cdot B2)$$

Ladder Logic Representation



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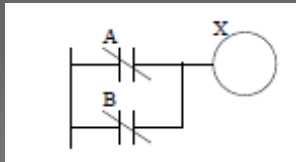
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Simplified Forms of Functions

Avoid multiple complemented variables in ladder logic (No NAND, NOR)

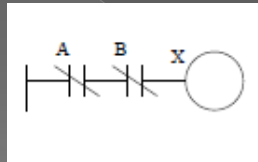
NAND

$$X = \overline{A \cdot B} = \overline{A} + \overline{B}$$



NOR

$$X = \overline{A + B} = \overline{A} \cdot \overline{B}$$



NAND/NOR can not be implemented effectively using software.
(Programmable Logic Controllers)

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End Lesson 15: Boolean Representation of Ladder Diagrams

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